

# The Classical Approach to Soergel Bimodules

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Setup:  $(W, S)$  Coxeter system

$\mathfrak{h}$  (geometric) realization of  $(W, S)$   
 $R = S(\mathfrak{h}^*) =$  polynomial functions on  $\mathfrak{h}$

$W \curvearrowright R$  &  $R^S = s\text{-invt polys}$

Consider all of  $R\text{-Bim}$  for the moment.

Standard bimodules:  $\nearrow$  graded  $R\text{-bimodules}$

Let  $x \in W$ .

Def: Define  $R_x \in R\text{-Bim}$  by

①  $R_x = R$  as a left  $R$ -module

② As a right  $R$ -module,  $m \cdot r = x(r)m$ .  
 $\uparrow \quad \uparrow$   
 $R_x \quad R$

Visualization:

Recall  $B_S$  can be viewed as

"border patrol"

$B_S$   
 $\downarrow$   
 semi-porous wall  
 where  $s\text{-invt polys}$   
 can slide through

For  $R_x$ , we have

"puberty"

$R_x$   
 $\downarrow$   
 $f = x(f)$   
 $\downarrow$   
 completely-porous wall  
 where anything can slide  
 through but act by  $x$

Facts:

①  $R_x \otimes_R R_y \cong R_{xy}$

②  $\text{Hom}(R_x, R_y) = \begin{cases} R & \text{if } x=y \\ 0 & \text{else} \end{cases}$

since  $\mathfrak{h}$  is a faithful rep

Define  $\text{StdBim} = \text{full } (\oplus, (n), \otimes)$  subcat of  $R\text{-Bim}$   
*additive*, *graded*, *tensor*  
 gen by  $\{R_x\}_{x \in W}$ .

① & ②  $\Rightarrow$  StdBim is a realization of the additive 2-groupoid of  $W$  over  $R$ .

- Tensor is multiplication, but

- Now we have more than just id maps, have mult by elts of  $R$  (polys in boxes)



with

$$\left| \begin{array}{c} \boxed{f} \\ \hline s \end{array} \right| = \left| \begin{array}{c} \boxed{sf} \\ \hline s \end{array} \right|$$

$\forall s \in S$ .

Filtrations:

Consider the SESs

$$\begin{array}{ccccc} 1 & \xrightarrow{\quad} & d_s := \frac{1}{2}(\alpha_s \otimes 1 - 1 \otimes \alpha_s) & & \\ (\Delta) & R_s(-1) & \xrightarrow{\quad} & B_s & \xrightarrow{\quad} & R_{id}(1) \\ & & & f \otimes g & \xrightarrow{\quad} & fg \end{array}$$

$$\begin{array}{ccccc} 1 & \xrightarrow{\quad} & c_s := \frac{1}{2}(\alpha_s \otimes 1 + 1 \otimes \alpha_s) & & \\ (\nabla) & R_{id}(-1) & \xrightarrow{\quad} & B_s & \xrightarrow{\quad} & R_s(1) \\ & & & f \otimes g & \xrightarrow{\quad} & fs(g) \end{array}$$

In alternate notation:  $f \otimes g \xrightarrow{\quad} f \otimes g \leftarrow$  here  $\otimes$  is the wall | from before

Defining props of  $c_s$  &  $d_s$

$$fc_s = c_s f$$

$$fd_s = d_s s(f) \quad \forall f \in R$$

So,  $B_S$  is filtered by both  $R_{id}$  and  $R_S$ ,  
 but with no particular order (& grading shifts  
 depend on this  
 order).

Soon, we will fix this by specifying an order.

Def: A bimodule  $M \in R\text{-Bim}$  is said to have  
 property  $(*)$  if it has a finite filtration with  
 subquotients  $\{\oplus R_x(\cdot)\}$ .

Facts:

① Having  $(*)$  is closed under  $\otimes$ .

Suppose  $0 \subseteq B^1 \subseteq \dots \subseteq B^m = B$  w/ subquots  $\{\oplus R_x(\cdot)\}$   
 &  $0 \subseteq C^1 \subseteq \dots \subseteq C^n = C$

Lemma: If  $M$  can be filtered so that subquots  
 have  $(*)$ , then  $M$  has  $(*)$ .

$\forall i, j, B^{i+1}/B^i \otimes C^{j+1}/C^j$  is a std bimod,

so  $B \otimes C^{j+1}/C^j$  has  $(*)$ .

By Lemma,  $B \otimes C$  has  $(*)$ .

② Having  $(*)$  is closed under direct summands.

We are working in a Krull-Schmidt cat.

A filt of  $B \oplus C$  by indec subquots

( $R_x$  is indec) will give a filt of  $B$  &  $C$

separately.

If more  
 explanation  
 is needed.

## Observations:

i) Bott-Samelsons have (\*) by SESs & ①

ii) Summands of Bott-Samelsons (& therefore by ②  
Soergel bimods) have (\*)

iii)  $R_x$  are not Soergel bimods.

They can appear as submods & quot of  
Bott-Samelsons (as in SESs from before),  
but not as summands,

(Except  $R = R_{id} = B_{id}$ ).

Time to pin down this order/grading shift issue!

## $\Delta$ - & $\nabla$ -filtrations:

Def: A  $\Delta$ -filtration <sup>(or standard)</sup> on  $B \in \mathcal{S}Bim$  is a finite  
filtration by  $R$ -Bim  $0 \subset B^m \subset \dots \subset B^1 \subset B^0 = B$   
such that

$$B^i / B^{i+1} \cong \bigoplus R_x(v) \text{ with } l(x) = i \text{ \& } v \in \mathbb{Z}.$$

Analogously,

Def: A  $\nabla$ -filtration <sup>(or costandard)</sup> on  $B \in \mathcal{S}Bim$  is a finite  
filtration by  $R$ -Bim  $0 \subset B^0 \subset \dots \subset B^m = B$   
such that

$$B^i / B^{i-1} \cong \bigoplus R_x(v) \text{ with } l(x) = i \text{ \& } v \in \mathbb{Z}.$$

NOTE: The SES  $(\Delta)$  gives a  $\Delta$ -filtration on  $B_s$  &

$(\nabla)$  gives a  $\nabla$ -filtration on  $B_s$ .

Thm (Soergel 2006)

Any  $B \in \mathcal{S}\text{Bim}$  has a unique  $\Delta/\nabla$ -filt.

PF (Sketchy Sketch):

Having  $(\Delta)$  is preserved under taking summands  
(same reason as for  $(*)$ ).

Having  $(\Delta)$  is closed under  $\otimes$  with  $B_s$ .

$\hookrightarrow$  Gives  $(\Delta)$  for all Bott-Samelsons.  $\square$

The subquots do not depend on the choice of refinement,  
but order & grading shifts are very different!

Remark: Support filtrations give an explicit construction of  
 $\Delta/\nabla$ -filtrations!

One can define two ~~notions~~ notions of character.  
We will use  $\text{ch}_\Delta$ .

$$\text{ch}_\Delta: \mathcal{S}\text{Bim} \rightarrow H$$

$$\text{ch}_\Delta(B) = \sum_{x \in W} v^{\ell(x) + \text{shift on } R_x} H_x$$

E.g.: From SES,  $\text{ch}_\Delta(B_s) = v \text{Id} + H_s = \underline{H}_s$ .

Remark: ① This gives an inverse to the isomorphism

$$\varepsilon: H \xrightarrow{\sim} [\mathcal{S}\text{Bim}]$$

$$\underline{H}_s \longmapsto [B_s]$$

② In the course of the proof of the theorem  
above, Soergel showed

$$\text{ch}_\Delta(M \otimes B_s) = \text{ch}_\Delta(M) \cdot \underline{H}_s$$

Recall: Soergel conj:  $\forall x \in W, \text{ch}_\Delta(B_x) = \underline{H}_x$ .

## Localization:

Soergel bimodules become much simpler after localization, (become "like" standards).

Let  $Q$  be the homogeneous fraction field of  $R$ .

Even though  $Q$  is graded,  $Q \cong Q(2)$  in this case.

Def: The "localization" of  $B \in R\text{-Bim}$  is  $B \otimes_R Q$ .

NOTATION:  $B_s^Q := B_s \otimes_R Q$ ,  $BS^Q(\underline{w}) := BS(\underline{w}) \otimes_R Q$ ,

$$Q_{s_i} := R_{s_i} \otimes_R Q.$$

NOTE THAT:  $BS^Q(\underline{w}) \cong Q \otimes_{Q^{s_1}} \cdots \otimes_{Q^{s_n}} Q(n)$  if  $\underline{w} = s_1 \cdots s_n$

Observation:

Suppose  $M$  is free as an  $R$ -module.

Then, it includes into its own localization  
(injection since  $M$  is free).

But  $\text{Hom}(B_x, B_y)$  is free as a right  $R$ -module,

so get injection:

$$\text{Hom}(B_x, B_y) \hookrightarrow \text{Hom}(B_x, B_y) \otimes_R Q.$$

Localization is a faithful functor!

Back to SESs:

After localization, these SESs split each other:

$$\begin{array}{ccccc} R_s(-1) & \longrightarrow & B_s & \longrightarrow & R_s(1) \\ \parallel & \longrightarrow & \alpha_s = \frac{1}{2}(\alpha_s \otimes 1 - 1 \otimes \alpha_s) & \longrightarrow & \alpha_s \end{array}$$

&

$$\begin{array}{ccccc} R_{id}(-1) & \longrightarrow & B_s & \longrightarrow & R_{id}(1) \\ \parallel & \longrightarrow & \alpha_s = \frac{1}{2}(\alpha_s \otimes 1 + 1 \otimes \alpha_s) & \longrightarrow & \alpha_s \end{array}$$

But  $\alpha_s$  is  
invertible now!

So,  
 $B_s^Q \cong Q(1) \oplus Q_s(1).$

By expansion:

$$BS^{\mathbb{Q}}(\underline{w}) \cong \bigoplus_{e \leq \underline{w}} \mathbb{Q}_{\underline{w}_e}(\ell(\underline{w}))$$

↑  
This is key! Bott-Samelsons are easy after localization.

Why is localization useful?

① Proposition: For a rex  $\underline{w}$ ,  $\exists!$  indec summand  $B_{\underline{w}}$  in  $BS(\underline{w})$  which contains  $Q_{\underline{w}}$  after localization, and  $B_{\underline{w}}$  is not a summand of any shorter sequence.

Pf: Look at the decomposition of  $BS^{\mathbb{Q}}(\underline{w})$  above. Since  $Q_{\underline{w}}$  is indec and appears exactly once, there is certainly a unique summand which contains it. Again by looking above, it is clear that  $Q_{\underline{w}}$  does not appear in the localization of any shorter sequence.  $\square$

NOTE: This does not prove that all the other summands in  $BS(\underline{w})$  do appear in some shorter sequence.

② Localization is a faithful functor, so can check if two morphisms of Soergel bimods are equal by checking after localization.

- Geordie writes computer programs exploiting this fact.

③ There are nice diagrammatics for playing with localization!